

# Predicting many properties of a quantum system from very few measurements

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Journal Club 2025/6/11

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Reference : [\*Nat. Phys.\* \*\*16\*\*, 1050–1057 \(2020\)](#)

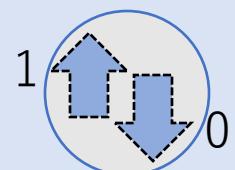
Authors : Hsin-Yuan Huang, Richard Kueng, and John Preskill

# 1. Introduction

## Quantum computer

Unit of information:

→ Qubit



$$\alpha|0\rangle + \beta|1\rangle$$



New computational methods that have been developed recently.



## Simulation of Quantum Physics

- Quantum many body system
- Lattice quantum field theory

When analytical calculations are not feasible, observables  $\langle \mathcal{O} \rangle$  are computed numerically

Energy, correlation function etc

Today's topic :

Efficiently extracting many physical observables from a quantum computer.

# 1. Introduction

Why do we need quantum computers?

## **The sign problem**

- In lattice QFT, people use the path-integral formalism and observables are typically evaluated using Monte Carlo integration
- At finite chemical potential or real-time dynamics, the path-integral weight becomes complex.
- As a result, important sampling breaks down → sign problem

# 1. Introduction

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## **The sign problem**

- In lattice QFT, people use the path-integral formalism and observables are typically evaluated using Monte Carlo integration
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- As a result, important sampling breaks down → sign problem

## **The limitations of tensor networks**

- To circumvent the sign problem, researchers are turning to the Hamiltonian formalism.
- Tensor network methods work well in (1+1)D systems where the entanglement follows an area law.
- However, they become inefficient in higher dimensions or in systems exhibiting volume-law entanglement, where the required bond dimension grows rapidly.

# 1. Introduction

Why do we need quantum computers?

**The sign problem**

**The limitations of tensor networks**



**Quantum computers can solve these problems.**

## Quantum computers

- Quantum computers are the computers built from quantum-mechanical components that obey the laws of quantum mechanics. **=Qubits**
- An  $n$ -qubit system can represent a  $2^n$ -dimensional Hilbert space directly.
- They are free from the sign problem.
- Highly entangled states can be prepared naturally.
- Real-time quantum dynamics is implemented directly as unitary time evolution.

# 1. Introduction

## Basics of quantum computer

### Quantum states

1-qubit state :  $|\psi\rangle = \alpha_1|0\rangle + \alpha_2|1\rangle$

Super position

2-qubits state :  $\alpha_{00}|00\rangle + \alpha_{01}|01\rangle + \alpha_{10}|10\rangle + \alpha_{11}|11\rangle$

⋮

### Basic quantum gate operators

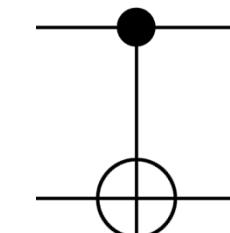
$$H = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}, X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \dots$$

$$CNOT = |0\rangle\langle 0| \otimes 1 + |1\rangle\langle 1| \otimes X, \dots$$

### Quantum circuit notation

$$|\psi\rangle \xrightarrow{U} U|\psi\rangle$$

$$U = H, X, Y \dots \text{etc}$$

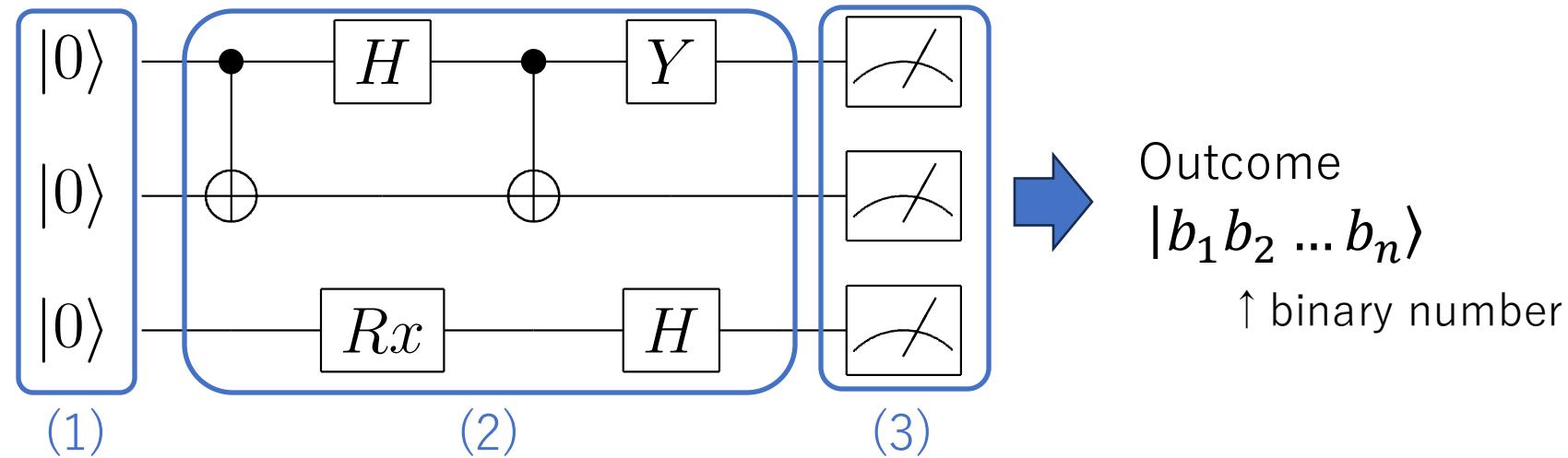


※Arbitrary unitary operators can be written in terms of basic quantum gates.

# 1. Introduction

Quantum computation consists of three essential steps:

- { (1) Prepare an initial state  $|0\rangle^{\otimes n}$
- (2) Apply a sequence of quantum gates
- (3) Measure in the computational basis

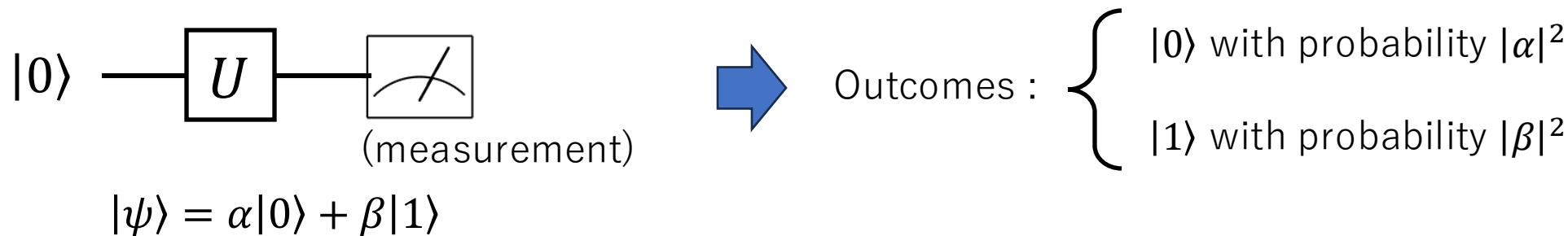


By repeatedly performing these three steps and collecting statistics from the measurement outcomes, we carry out quantum computation.

# 1. Introduction

The simple example : calculation of  $\langle Z \rangle$

Let us consider estimating expectation value  $\langle Z \rangle$  for the quantum state  $|\psi\rangle = U|0\rangle$ .



Repeating this experiment many times, we can estimate  $|\alpha|^2$  and  $|\beta|^2$  by

$$|\alpha|^2 \approx \frac{N_0}{N_{\text{shots}}}, \quad |\beta|^2 \approx \frac{N_1}{N_{\text{shots}}}$$

$N_{\text{shots}}$  : the total number of experiments  
 $N_{0/1}$  : the number of measurement outcome 0/1

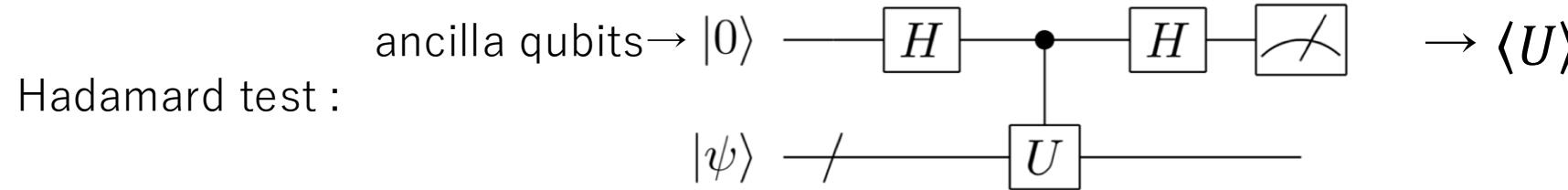
Then, the expectation value  $\langle Z \rangle$  is estimated as follows;

$$\langle Z \rangle = |\alpha|^2 - |\beta|^2 \approx \frac{N_0 - N_1}{N_{\text{shots}}}$$

# 1. Introduction

When we extend to multi-qubits systems, the computational possibilities expand dramatically.

By introducing ancilla qubits (auxiliary qubits), we can implement more sophisticated operations.



## Examples of quantum algorithms:

- Hadamard test : estimation of  $\langle U \rangle$  for arbitrary unitary  $U$
- SWAP test : estimation of  $\langle \psi | \phi \rangle$
- Quantum Fourier transform : discrete Fourier transformation
- Phase estimation : estimation of eigenvalue of unitary operator
- Grover's algorithm : search algorithm for solution  $f(x) = 1$

These quantum algorithms enable us to estimate relevant quantities and address specific computational tasks.

# 1. Introduction

Let me consider the following goal.

**Goal** : Estimating  $M$  different physical observables, where  $M$  is large.

$$\langle \mathcal{O}_i \rangle = \text{Tr}[\mathcal{O}_i \rho], \quad i = 1, 2, \dots, M \gg 1$$

We can estimate individual  $\langle \mathcal{O}_i \rangle$  by known quantum algorithm (e.g. Hadamard test)

However...

## Problems

- Naively, we need to design and execute  $M$  different quantum circuits.  
→△We must perform  $\mathcal{O}(M)$  measurements.
- Many algorithms also require ancilla qubits, increasing the total number of qubits and circuit depth.  
→△Current quantum computers are still small and contain noise.
- The required gate are often complex and difficult to implement on near term quantum computers.

# 1. Introduction

**Quantum state tomography** : reconstructing the full density matrix  $\rho$  of a quantum system.  
[Sugiyama-Turner-Murao 2013]

$$\rho \left\{ \begin{array}{c} \xrightarrow{\quad} X, Y, Z \\ \xrightarrow{\quad} X, Y, Z \\ \vdots \\ \xrightarrow{\quad} X, Y, Z \end{array} \right. \quad \xrightarrow{\quad} \quad \rho = \frac{1}{2^n} \sum_{i=1}^{4^n} \langle P_i \rangle P_i$$
$$P_i \in \{I, X, Y, Z\}^n$$

## Problems

- Needs to measure in many different basis (Pauli X, Y, Z)  
→△Required the number of measurements is  $\mathcal{O}(3^n)$  for  $n$ -qubits system.
- Postprocessing is computationally expensive ( $= \mathcal{O}(4^n)$ )  
→△Impractical for large system.



Quantum state tomography is Not efficient for estimating many observables

# 1. Introduction

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$$\langle \mathcal{O}_i \rangle = \text{Tr}[\mathcal{O}_i \rho], \quad i = 1, 2, \dots, M \gg 1$$

## Estimating $\langle \mathcal{O}_i \rangle$ by known quantum algorithm

- $M$  different quantum circuits
- # of measurement  $\sim \mathcal{O}(M)$
- Ancilla qubits
- Complex gate operation

  challenging

## Quantum state tomography

We can estimate  $\rho$  directly but,

- # of measurement  $\sim \mathcal{O}(3^n)$
- Classical postprocessing cost  $\sim \mathcal{O}(4^n)$

  challenging  
(Not scalable)

Are there efficient ways to estimate many observables from a quantum system?

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## Quantum state tomography

We can estimate  $\rho$  directly but,

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→  challenging  
(Not scalable)

Are there efficient ways to estimate many observables from a quantum system?

→ **Classical shadow** [Huang-Kueng-Preskill 2020]

# 1. Introduction

## Short summary of classical shadow

- Estimates  $M$  different physical observables from only  $\mathcal{O}(\log M)$  measurements.  
→ **Exponential speed up**
- Does not use ancilla qubits or multiple copies of system.  
→ **The number of required qubits is minimal.**
- Requires only shallow randomized circuits  
→ **Suitable for current quantum devices**
- The authors demonstrate its effectiveness across various observables, including two-point correlation functions, energy, and entanglement entropy.

# Outline

1. Introduction

2. Classical shadow

3. Applications and numerical results

4. Summary

# Outline

1. Introduction

2. Classical shadow

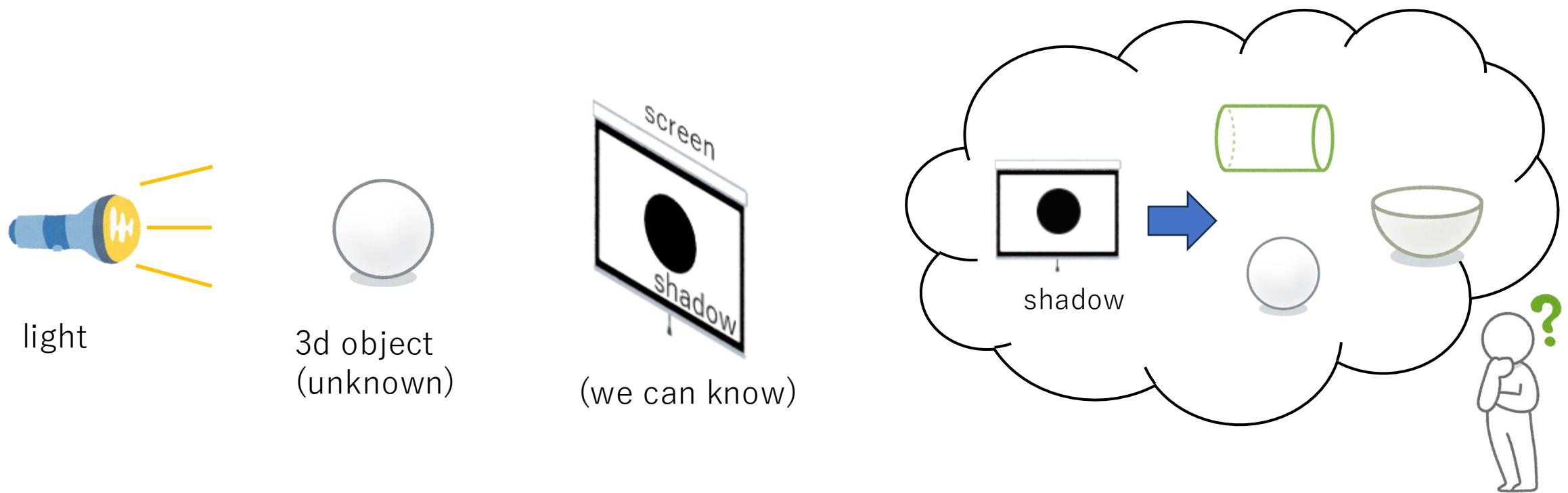
3. Applications and numerical results

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## 2. Classical shadow

Before introducing the algorithm of classical shadow, let me explain a intuition.

Reconstructing a 3d object from 2d shadows.

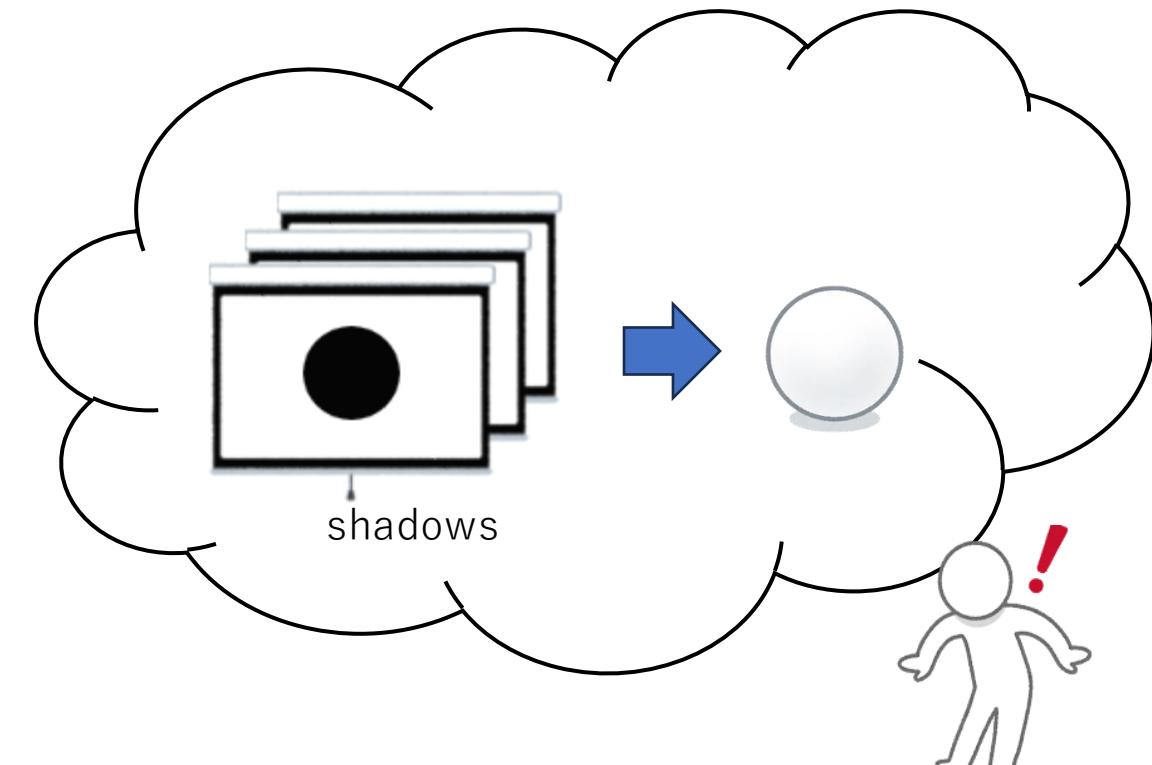
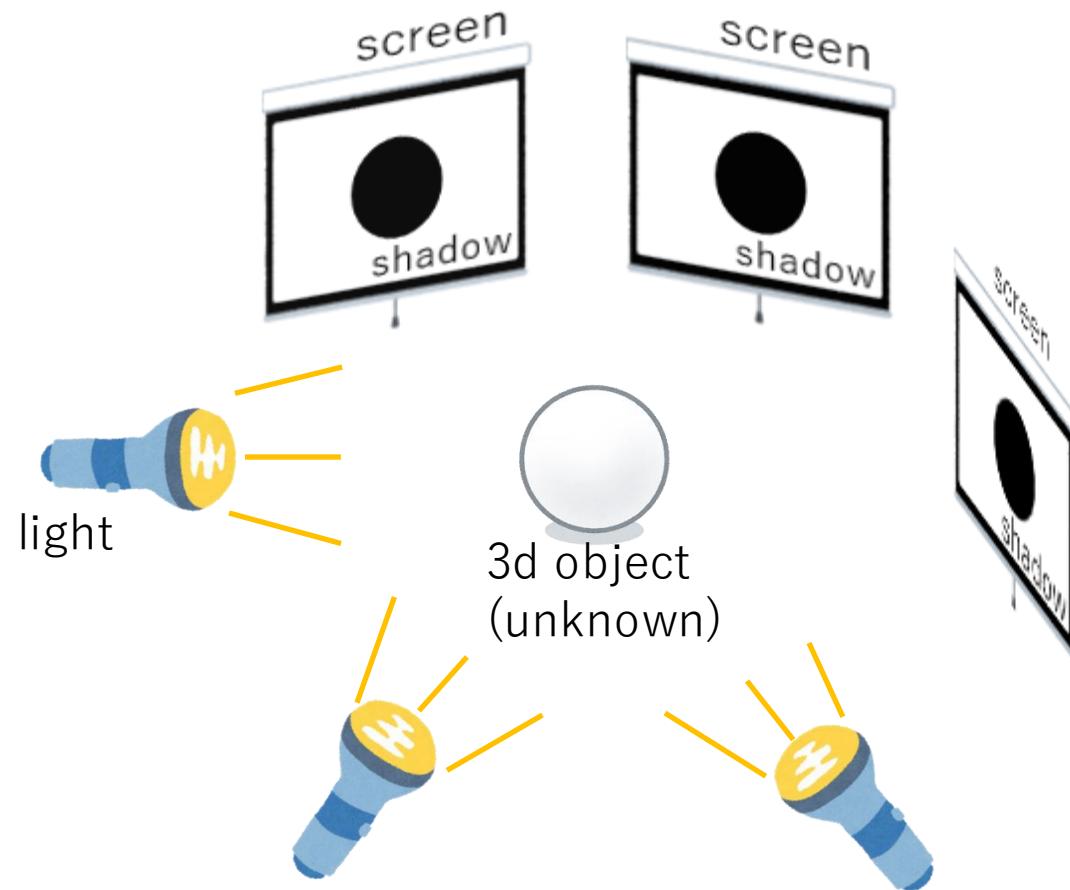


**We cannot reconstruct the 3d object from single shadow.**

## 2. Classical shadow

Before introducing the algorithm of classical shadow, we get intuition.

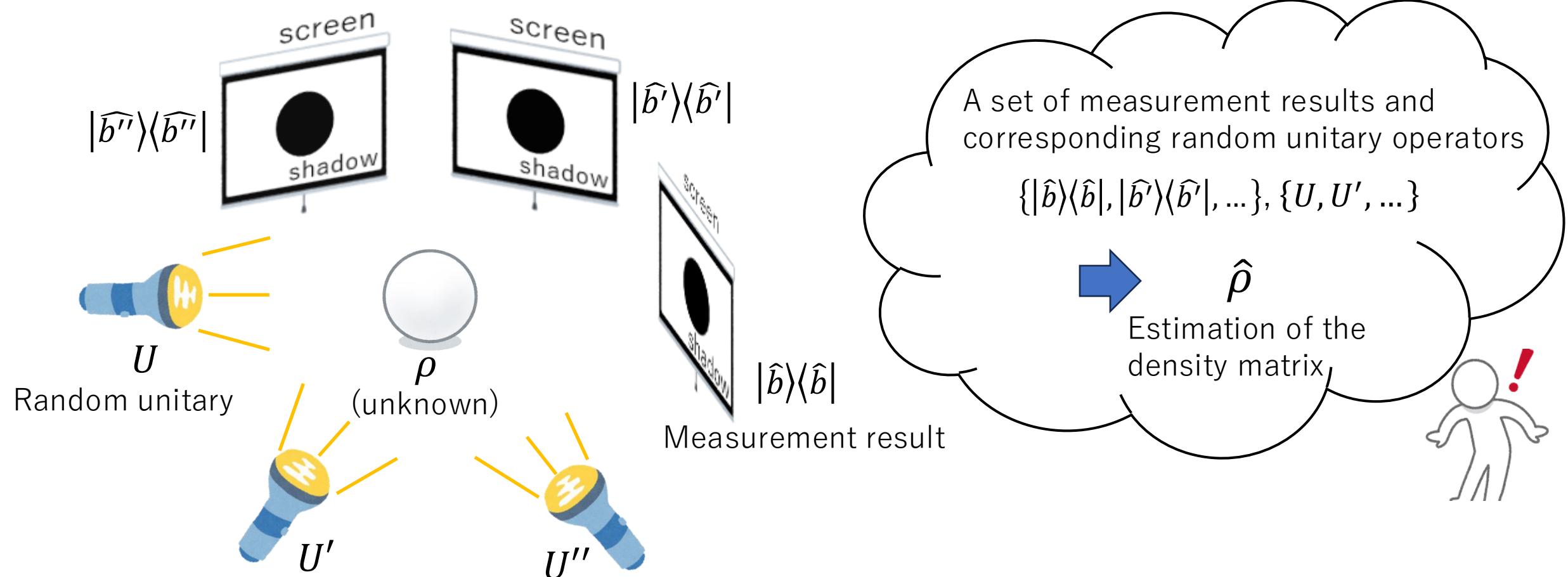
Reconstructing a 3d object from 2d shadows.



**If we collect many shadows from various directions, we can estimate the 3d object.**

## 2. Classical shadow

The idea of classical shadow:

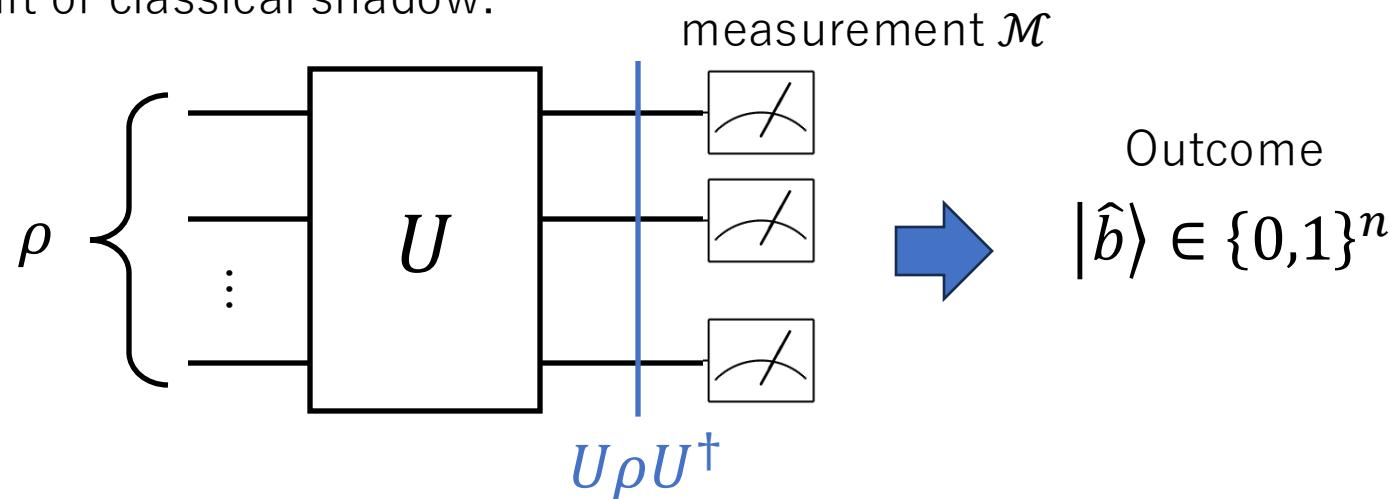


If we collect many measurements results with various basis, we can estimate the density matrix.

## 2. Classical shadow

Consider the  $n$  qubits system.

The quantum circuit of classical shadow:



### Procedures

1. Choose a random unitary  $U \in \mathcal{U}$  and apply to the state : $\rho \mapsto U\rho U^\dagger$
2. Perform a computational basis measurement  $\mathcal{M}$  and get outcome  $|\hat{b}\rangle$
3. Restore the classical snapshot  $U^\dagger|\hat{b}\rangle\langle\hat{b}|U$

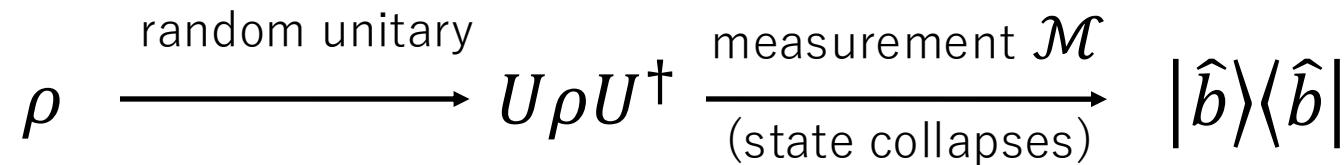
Repeating this procedures  $N$  times,

$$\mathbb{E}_{u,b}[U^\dagger|\hat{b}\rangle\langle\hat{b}|U] = \mathcal{M}(\rho)$$

Ensemble average over  $\mathcal{U}$  and outcome  $b \in \{0,1\}^n$

True density matrix

## 2. Classical shadow



### The idea of classical shadow

$$\text{Classical shadow : } \hat{\rho} = \mathcal{M}^{-1}(U^\dagger |\hat{b}\rangle\langle\hat{b}| U)$$

“A classical guess of the density matrix based on a single measurement outcome”  
(where  $\mathcal{M}^{-1}$  depends on the unitary ensemble  $\mathcal{U}$  )



$$\rho = \mathbb{E}_{U,b}[\hat{\rho}]$$

The ensemble average gives the true density matrix

**Classical shadow is a classical approximation of the density matrix constructed from a measurement outcome.**

## 2. Classical shadow

Example of 1-qubit classical shadow  $\hat{\rho}$

For  $n = 1$  case, we can choose  $\mathcal{U}$  as Clifford random unitaries  $\text{Cl}(2)$  :

$$\text{Cl}(2) = \{U \in \text{U}(2) \mid \forall P \in \mathcal{P}, UPU^\dagger \in \mathcal{P}\} \text{, where } \mathcal{P} \text{ is Pauli group}$$

Classical shadow per single measurement:

$$\hat{\rho} = \mathcal{M}^{-1}(U^\dagger |\hat{b}\rangle\langle \hat{b}|U) = 3U^\dagger |\hat{b}\rangle\langle \hat{b}|U - I \quad , \text{ where } U \in \text{Cl}(2)$$

It is mathematically shown that this classical shadow reproduces the density matrix  $\rho$ .

$$\mathbb{E}_{U,b}[\hat{\rho}] = \rho$$

※ This result follows from the fact that Clifford unitaries form a unitary 2-design.

$$\mathbb{E}_{U,b}[\hat{\rho}] = \mathbb{E}_U \left[ \sum_b \underbrace{\langle b|U\rho U^\dagger|b\rangle}_{\text{probability}} (3U^\dagger|b\rangle\langle b|U - I) \right] = \rho$$

Unitary 2-design

## 2. Classical shadow

Based on this idea, the following algorithm was proposed: [Huang-Kueng-Preskill 2020]

### Algorithm

1. Perform the experiment  $N$  times and collect the classical shadows:

$$\{\hat{\rho}_1, \hat{\rho}_2, \dots, \hat{\rho}_N\}, \text{ where } \hat{\rho}_k = \mathcal{M}^{-1}(U_k^\dagger |\hat{b}_k\rangle \langle \hat{b}_k| U_k), \text{ for } k = 1, \dots, N$$

2. Estimate the expectation values of  $M$  observables using a classical computer :

$$\langle \hat{O}_i \rangle = \frac{1}{N} \sum_{k=1}^N \text{Tr}[\mathcal{O}_i \hat{\rho}_k], \text{ for } i = 1, \dots, M$$

3. Increase the number of experiments  $N$  until the statistical error reaches the desired accuracy  $\epsilon$ .

※ The required classical memory is  $\mathcal{O}(NM)$

**A finite number of classical shadows is sufficient to estimate observables.**

## 2. Classical shadow

Theoretical guarantee of Classical shadows:

**Theorem 1** [Huang-Kueng-Preskill 2020]

To estimate  $M$  observables up to accuracy  $\epsilon$ , the required number of measurements  $N$  is given by:

$$N = \mathcal{O}\left(\frac{\log(M)}{\epsilon^2} \max_i \|O_i\|_{\text{shadow}}\right)$$

where  $\|O_i\|_{\text{shadow}}$  depends on the unitary ensemble  $\mathcal{U}$  used to generate the classical shadows.

Examples of shadow norms for different unitary ensembles:

$$\begin{cases} \mathcal{U} = \text{Cl}(2^n) \Rightarrow \|O_i\|_{\text{shadow}}^2 \leq \text{Tr}[O_i^2] \\ \mathcal{U} = \text{Cl}(2)^{\otimes n} \Rightarrow \|O_i\|_{\text{shadow}}^2 \leq 4^k \|O_i\|^2 \end{cases} \quad \|O_i\| : \text{operator norm}$$

where  $k$  denote the locality of  $O_i$  , e.g.  $O_i = X_1 \otimes Y_2 \Rightarrow k = 2$ .

- The number of required measurements **scales only logarithmically** with the number of observables  $M$ .
- The shadow norm determines the difficulty of estimating each observable.

# Outline

1. Introduction

2. Classical shadow

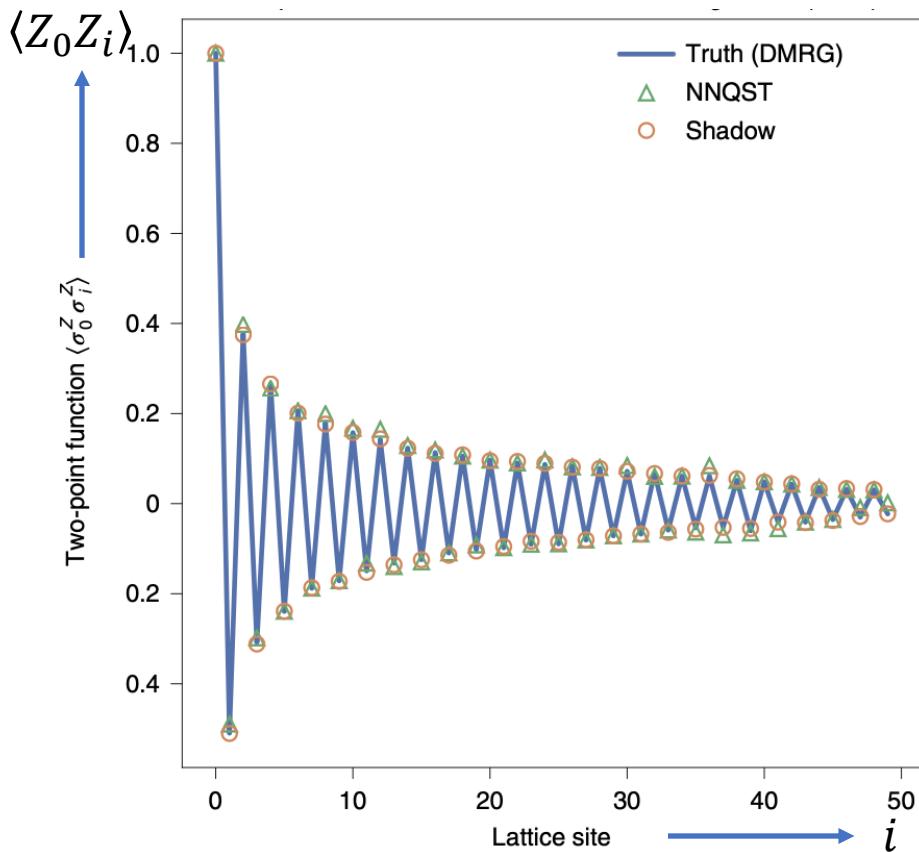
3. Applications and numerical results

4. Summary

### 3. Applications and numerical results

#### Application 1 : 1d transverse Ising model

$$H = J \sum_j Z_j Z_{j+1} + h \sum_j X_j \quad \text{, where we set } J = h.$$



**Goal** : two-point functions  $\langle Z_0 Z_i \rangle$  for  $i = 1, \dots, N_{\text{site}} = 50$

— Exact results(DMRG)

△ Quantum state tomography [Torlai, et al 2018]

○ Classical shadow

We use  $\mathcal{U} = \text{Cl}(2)^{\otimes N_{\text{site}}}$  and  $N = 2^{19}$  measurement snapshots.

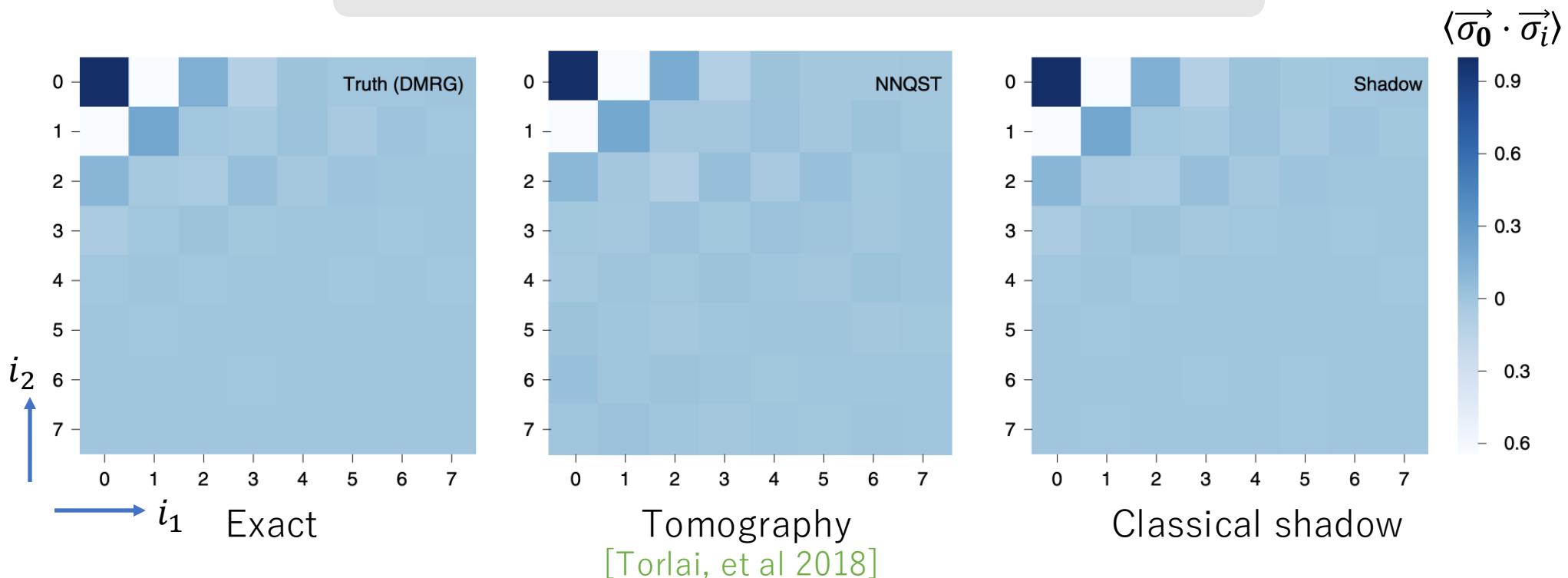
The classical shadow predictions perfectly match the exact results.

### 3. Applications and numerical results

## Application 2: 2d Heisenberg model

$$H = J \sum_{\langle i,j \rangle} \vec{\sigma}_i \cdot \vec{\sigma}_j \quad \text{with an } 8 \times 8 \text{ triangular lattice}$$

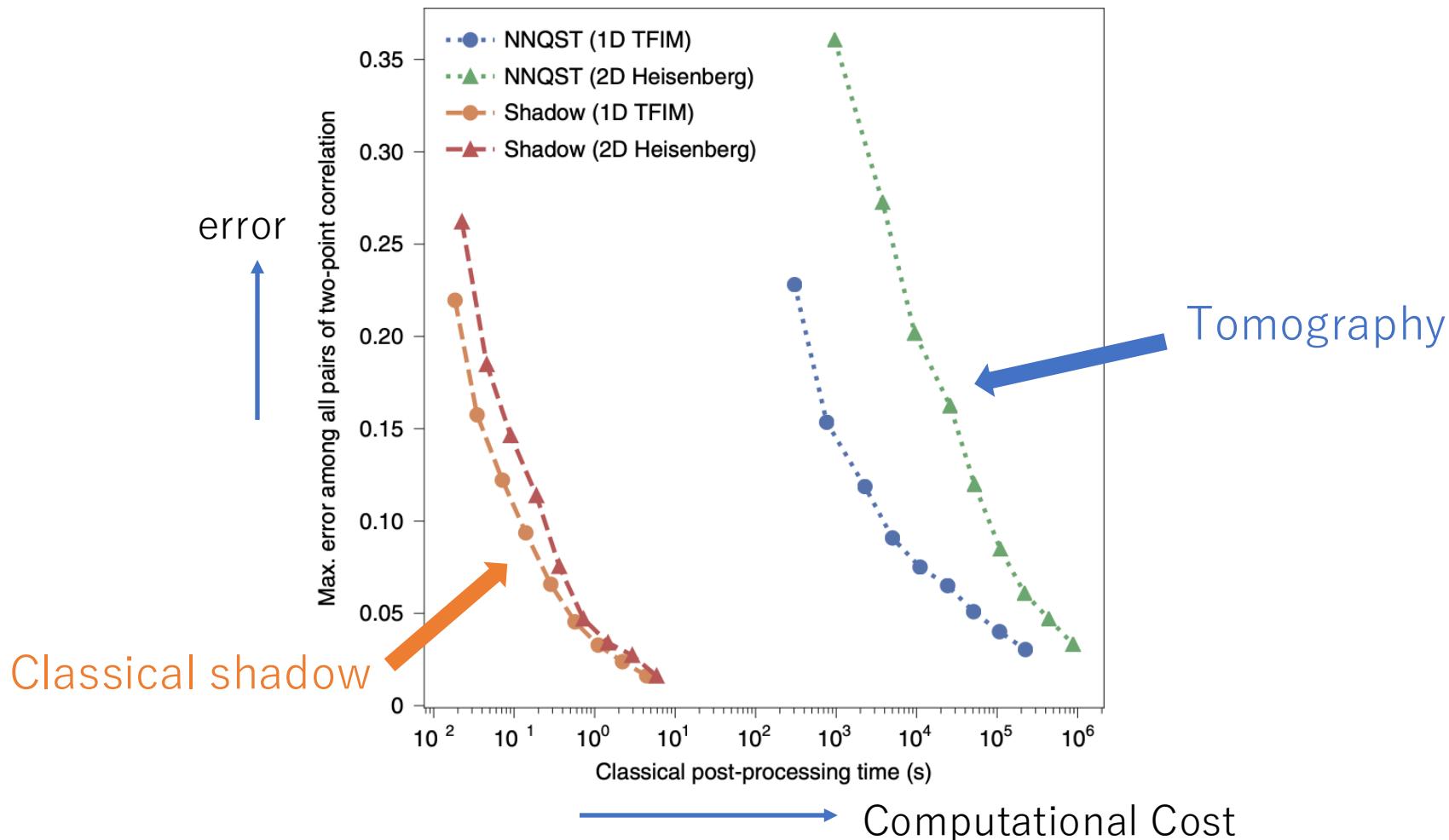
**Goal:** two-point functions  $\langle \vec{\sigma}_0 \cdot \vec{\sigma}_i \rangle$  for  $i = (i_1, i_2)$



The classical shadow predictions perfectly match the exact results.

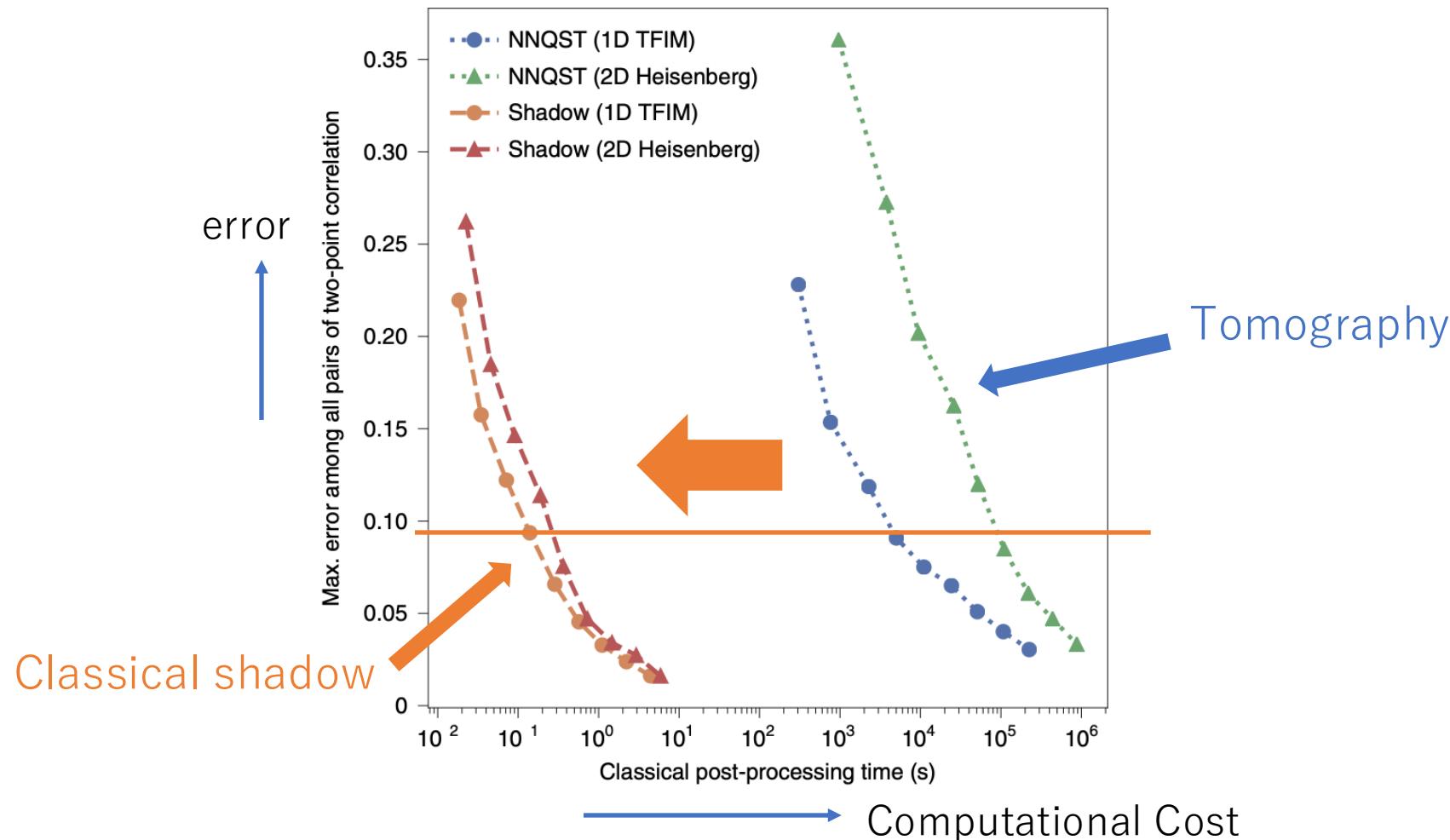
### 3. Applications and numerical results

#### Computational cost of application 1 and 2



### 3. Applications and numerical results

#### Computational cost of application 1 and 2



Classical shadows achieve comparable accuracy with significantly fewer measurements.

### 3. Applications and numerical results

#### Application 3: Ground state energy estimation in the Schwinger model

Jordan-Wigner transformation

$$\mathcal{L} = -\frac{1}{4} F_{\mu\nu} F^{\mu\nu} + i\bar{\psi} \gamma^\mu (\partial_\mu + ig A_\mu) \psi - m\bar{\psi} \psi \quad \longrightarrow \quad H_{\text{spin}} : \text{spin Hamiltonian}$$

**Goal:** ground state energy

The ground state of the Schwinger model was studied by VQE in [Kokail, C. et al. 2019].

#### VQE

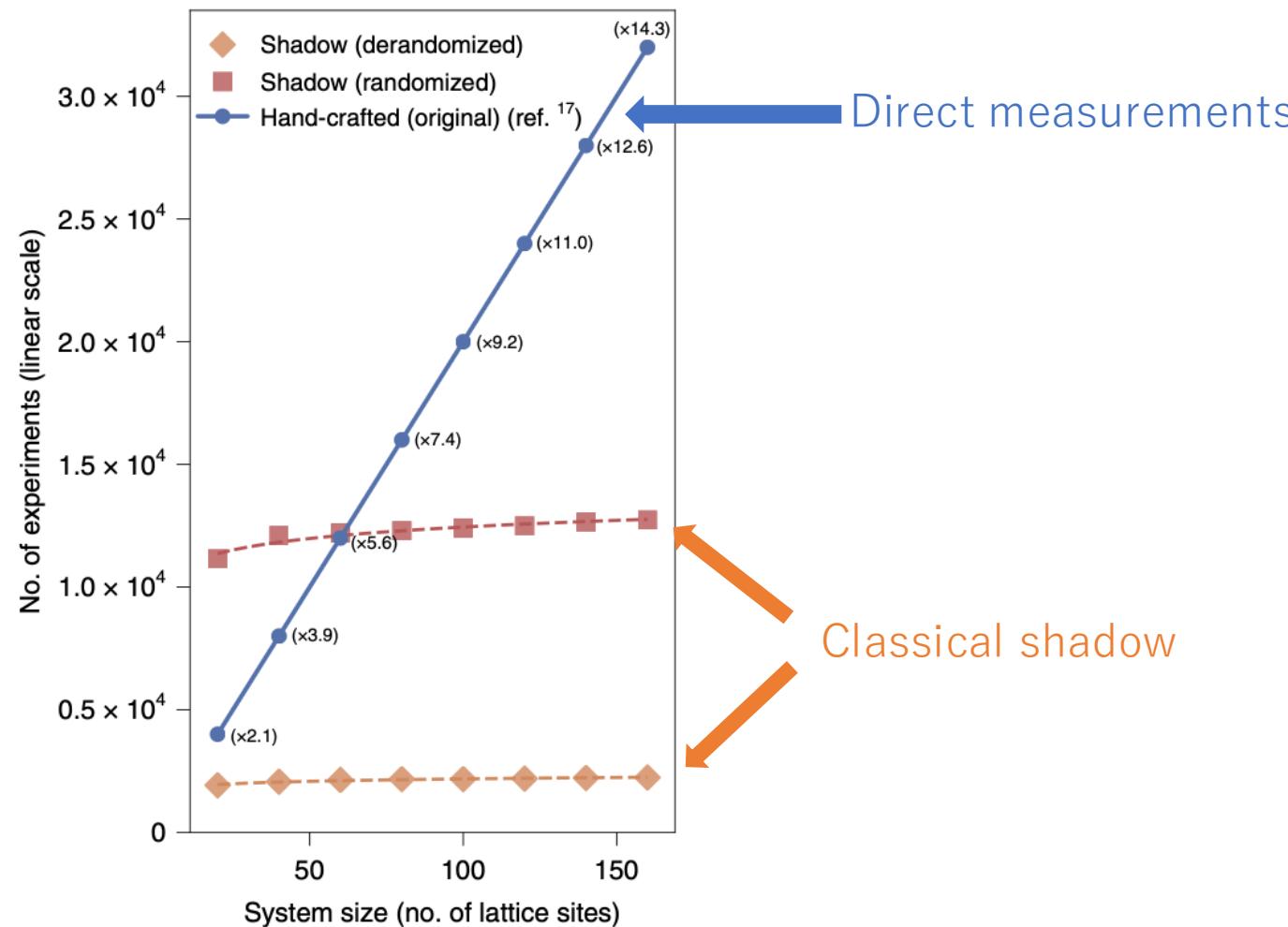
1. Set ansatz  $|\psi(\theta)\rangle$ ,  $\theta$ : parameters
2. Calculate  $\langle H_{\text{spin}} \rangle_\theta$  by a quantum computer
3. Update parameters  $\theta$  to minimize the energy  $\langle H \rangle_\theta$

We compare classical shadow-based energy estimation with the standard method using direct measurements as implemented in [Kokail, C. et al. 2019].

### 3. Applications and numerical results

#### Application 3: Ground state energy estimation in the Schwinger model

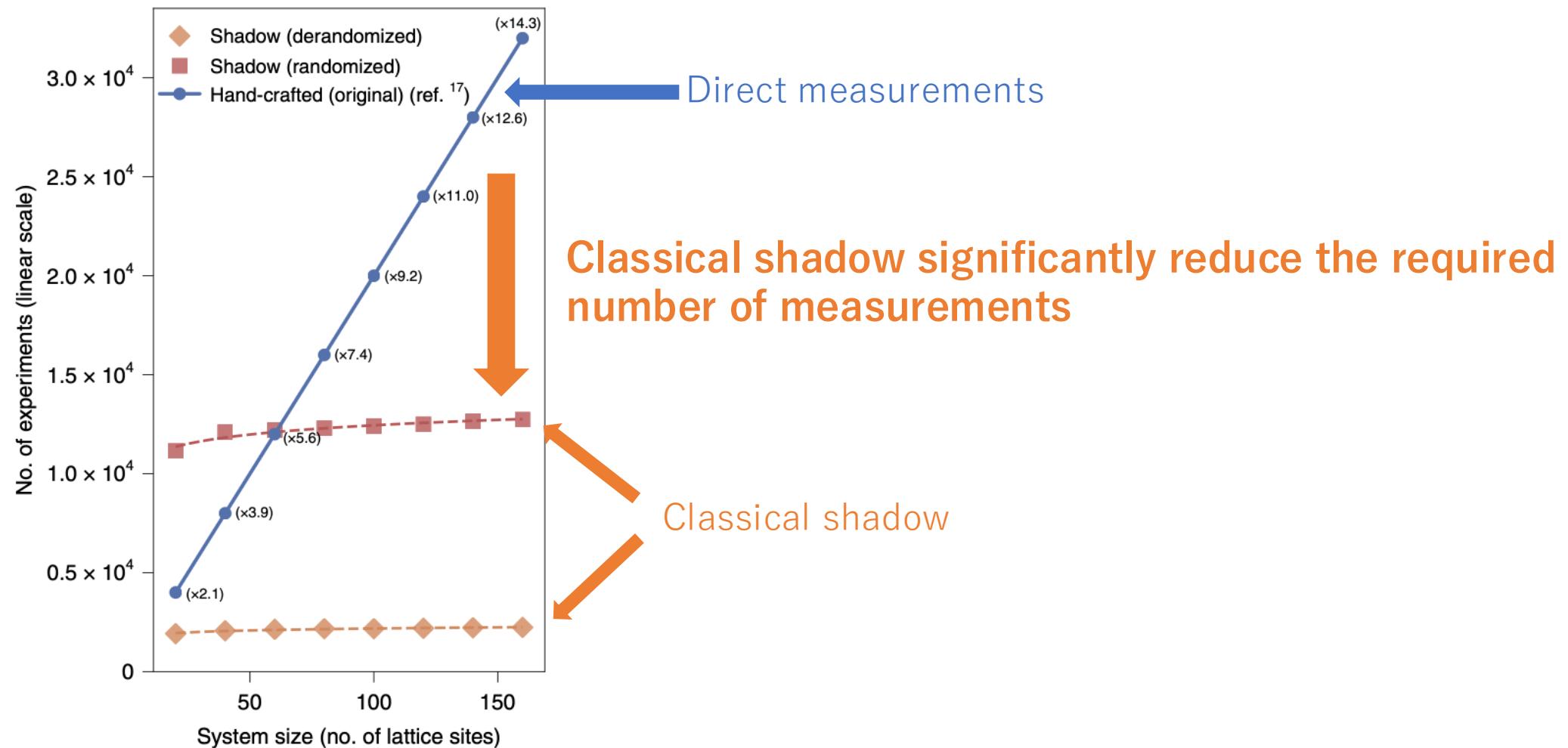
Scaling of the number of measurements required to achieve a fixed error.



### 3. Applications and numerical results

#### Application 3: Ground state energy estimation in the Schwinger model

Scaling of the number of measurements required to achieve a fixed error.



### 3. Applications and numerical results

#### Application 4: Estimating entanglement Rényi entropy

Using classical shadows, one can estimate the purity  $\text{Tr}[\rho_A^2]$ , which is related to the second Rényi entanglement entropy.

##### Estimation of purity by classical shadows

$$\text{Tr}[\rho_A^2] \approx \frac{1}{N(N-1)} \sum_{i \neq j}^N \text{Tr}[\hat{\rho}_{A,i} \hat{\rho}_{A,j}]$$

$\{\hat{\rho}_{A,1}, \dots, \hat{\rho}_{A,N}\}$ : classical shadows on subsystem  $A$ .

Calculated via pairwise trace overlaps between shadow estimates.

The number of required measurements:

$$N = \mathcal{O}\left(\frac{2^n}{\epsilon^2}\right)$$

$n$ : number of qubits  
 $\epsilon$ : desired statistical error

This cost arises from the non-local nature of entanglement, and the authors showed that it is unavoidable due to fundamental information-theoretical bounds.

# Outline

1. Introduction

2. Quantum Algorithm : Classical shadow

3. Applications and numerical results

4. Summary

## 4. Summary

- Classical shadow efficiently estimates many observables from quantum states.
- Only  $\mathcal{O}(\log M)$  measurements are needed to estimate  $M$  observables  $\langle \mathcal{O}_i \rangle = \text{Tr}[\mathcal{O}_i \rho]$ .  
→ **Exponential speed up!**  $i = 1, 2, \dots, M$
- Requires only shallow quantum circuits → **Suitable for near-term quantum computer**
- Estimating non-local properties such as entanglement requires many measurements, but it is efficient among methods that do not use ancilla or additional qubits.
- **Classical shadows are promising tools for studying quantum many-body systems and quantum field theory simulations on near-term quantum computers.**

# Appendix

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## The concrete form of classical shadows

$$u = \text{Cl}(2^n) \Rightarrow \hat{\rho} = (2^n + 1)U^\dagger |\hat{b}\rangle\langle\hat{b}|U - \mathbb{I}$$

$$u = \text{Cl}(2)^{\otimes n} \Rightarrow \hat{\rho} = \bigotimes_{j=1}^n \left( 3U_j^\dagger |\hat{b}_j\rangle\langle\hat{b}_j|U_j - \mathbb{I} \right)$$

# Appendix

Random unitary ensemble  $\mathcal{U}$  should be tomographically complete.

For  $\rho \neq \sigma$ , there exist  $U \in \mathcal{U}$  and  $b$  s.t.

$$\langle b | U\rho U^\dagger | b \rangle \neq \langle b | U\sigma U^\dagger | b \rangle$$